## Optimization Methods - Duality

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## 1 Solutions

**Problem 1:** Get the Lagrange function and dual function of the following problem:

$$\min x^T x 
\text{s.t. } Ax = b$$
(1)

**Solution:** This problem has no inequality constraints and p (linear) equality constraints. The Lagrangian is  $L(x,\nu) = x^T x + \nu^T (Ax - b)$ , with domain  $\mathbb{R}^n \times \mathbb{R}^p$ . The dual function is given by  $g(\nu) = \inf_x L(x,\nu)$ . Since  $L(x,\nu)$  is a convex quadratic function of x, we can find the minimizing x from the optimality condition

$$\nabla_x L(x, \nu) = 2x + A^T \nu = 0,$$

which yields  $x = -(1/2)A^T\nu$ . Therefore, the dual function is

$$g(\nu) = L(-(1/2)A^T\nu, \nu) = -(1/4)\nu^T AA^T\nu - b^T\nu,$$

which is a concave quadratic function, with domain  $\mathbb{R}^p$ . The lower bound property states that for any  $\nu \in \mathbb{R}^p$ , we have

$$-(1/4)\nu^{T}AA^{T}\nu - b^{T}\nu \le \inf\{x^{T}x|Ax = b\}$$

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**Problem 2:** Get the Lagrange function and dual function of the following problem:

$$\min c^T x$$
s.t.  $Ax = b, x > 0$ . (2)

**Solution:** To form the Lagrangian we introduce multipliers  $\lambda_i$  for the *n* inequality constraints and multiplier  $\nu_i$  for the equality constraints, and obtain

$$L(x, \lambda, \nu) = c^{T}x - \sum_{i=1}^{n} \lambda_{i}x_{i} + \nu^{T}(Ax - b) = -b^{T}\nu + (c + A^{T}\nu - \lambda)^{T}x$$

The dual function is

$$g(\lambda, \nu) = \inf L(x, \lambda, \nu) = -b^T \nu + \inf (c + A^T \nu - \lambda)^T$$

which is easily determined analytically, since a linear function is bounded below only when it is identically zero. Thus,  $g(\lambda, \nu) = -\infty$  except when  $c + A^T \nu - \lambda = 0$ , in which case it is  $-b^T \nu$ :

$$g(\lambda, \nu) = \begin{cases} -b^T \nu & A^T \nu - \lambda + c = 0, \\ -\infty & otherwise. \end{cases}$$
 (3)

Note that the dual function g is finite only on a proper affine subset of  $\mathbb{R}^m \times \mathbb{R}^p$ . We will see that this is a common occurrence. The lower bound property (5.2) is nontrivial only when  $\lambda$  and  $\nu$  satisfy  $\lambda \geq 0$  and  $A^T \nu - \lambda + c = 0$ . When this occurs,  $-b^T \nu$  is a lower bound on the optimal value of the LP.

**Problem 3:** Get the optimal value of the following problem via KKT conditions:

$$\min (1/2)x^T P x + q^T x + r$$
s.t.  $Ax = b, P \in S_+^n$  (4)

**Solution:** The KKT conditions for this problem are

$$Ax^* = b, Px^* + q + A^T \nu^* = 0,$$

which we can write as

$$\left[\begin{array}{cc} P & A^T \\ A & 0 \end{array}\right] \left[\begin{array}{c} x^* \\ \nu^* \end{array}\right] = \left[\begin{array}{c} -q \\ b \end{array}\right]$$

Solving this set of m+n equations in the m+n variables  $x^*, \nu^*$  gives the optimal primal and dual variables.

**Problem 4:** Get the optimal value of the max entropy problem via KKT conditions:

$$\min \sum_{i=1}^{n} x_i \log x_i$$
s.t.  $Ax \le b, \mathbf{1}^T x = 1$  (5)

**Solution:** The Lagrange function is

$$L(x, \lambda, \nu) = \sum_{i=1}^{n} x_i \log x_i + \lambda^{T} (Ax - b) + \nu (\mathbf{1}^{T} x - 1)$$

By taking the derivative and setting it to zero:

$$\frac{\partial L}{\partial x_i} = \log x_i + 1 + a_i^T \lambda + \nu = 0$$

Therefore,  $x_i = e^{-1-a_i^T \lambda - \nu}$ . Substitute  $x_i$  in  $g(\lambda, \nu)$ 

$$g(\lambda, \nu) = \sum_{i=1}^{n} (x_i(-\lambda^T a_i - \nu - 1) + \lambda^T (a_i^T x_i - b_i) + \nu x_i) - \nu$$

$$= -b^T \lambda - \nu - e^{-\nu - 1} \sum_{i=1}^{n} e^{-a_i^T \lambda}$$
(6)

The dual problem is

$$\max -b^{T}\lambda - \nu - e^{-\nu - 1} \sum_{i=1}^{n} e^{-a_{i}^{T}\lambda}$$
s.t.  $\lambda > 0$  (7)

where  $a_i$  are the columns of A. We assume that the weak form of Slater's condition holds, i.e., there exists an x > 0 with  $Ax \le b$  and  $\mathbf{1}^T x = 1$ , so strong duality holds and an optimal solution  $(\lambda^*, \nu^*)$  exists. Suppose we have solved the dual problem. The Lagrangian at  $(\lambda^*, \nu^*)$  is

$$L(x, \lambda^*, \nu^*) = \sum_{i=1}^{n} x_i \log x_i + \lambda^{*T} (Ax - b) + \nu^* (\mathbf{1}^T x - 1)$$

which is strictly convex on D and bounded below, so it has a unique solution  $x^*$ , given by

$$x_i^* = 1/\exp(a_i^T \lambda^* + \nu^* + 1), i = 1, ..., n.$$

If  $x^*$  is primal feasible, it must be the optimal solution of the primal problem. If  $x^*$  is not primal feasible, then we can conclude that the primal optimum is not attained.

**Problem 5:** Get the optimal value of the water filling problem via KKT conditions:

$$\min -\sum_{i=1}^{n} \log(\alpha_i + x_i)$$
s.t.  $x > 0, \mathbf{1}^T x = 1$  (8)

**Solution:** This problem arises in information theory, in allocating power to a set of n communication channels. The variable  $x_i$  represents the transmitter power allocated to the *i*th channel, and  $\log(\alpha_i + x_i)$  gives the capacity or communication rate of the channel, so the problem is to allocate a total power of one to the channels, in order to maximize the total communication rate.

Introducing Lagrange multipliers  $\lambda^* \in \mathbb{R}^n$  for the inequality constraints  $x^* \geq 0$ , and a multiplier  $\nu^* \in \mathbb{R}$  for the equality constraint  $\mathbf{1}^T x = 1$ , we obtain the KKT conditions

$$x^* \ge 0, \mathbf{1}^T x^* = 1, \lambda^* \ge 0, \lambda_i^* x_i^* = 0, i = 1, ..., n, -1/(\alpha_i + x_i^*) - \lambda_i^* + \nu^* = 0, i = 1, ..., n.$$

We can directly solve these equations to find  $x^*, \lambda^*$ , and  $\nu^*$ . We start by noting that  $\lambda^*$  acts as a slack variable in the last equation, so it can be eliminated, leaving

$$x^* \ge 0, \mathbf{1}^T x^* = 1, x_i^* (\nu^* - 1/(\alpha_i + x_i^*)) = 0, i = 1, ..., n, \nu^* \ge 1/(\alpha_i + x_i^*), i = 1, ..., n.$$

If  $\nu^* < 1/\alpha_i$ , this last condition can only hold if  $x_i^* > 0$ , which by the third condition implies that  $\nu^* = 1/(\alpha_i + x_i^*)$ . Solving for  $x_i^*$ , we conclude that  $x_i^* = 1/\nu^* - \alpha_i$  if  $\nu^* < 1/\alpha i$ . If  $\nu^* \ge 1/\alpha_i$ , then  $x_i^* > 0$  is impossible, because it would imply  $\nu^* \ge 1/\alpha_i > 1/(\alpha_i + x_i^*)$ , which violates the complementary slackness condition. Therefore,  $x_i^* = 0$  if  $\nu^* \ge 1/\alpha_i$ . Thus we have

$$x_i^* = \begin{cases} 1/\nu^* - \alpha_i & \nu^* < 1/\alpha_i, \\ 0 & otherwise. \end{cases}$$
 (9)

or, put more simply,  $x_i^* = \max\{0, 1/\nu^* - \alpha_i\}$ . Substituting this expression for  $x_i^*$  into the condition  $\mathbf{1}^T x^* = 1$  we obtain

$$\sum_{i=1}^{n} \max\{0, 1/\nu^* - \alpha_i\} = 1.$$

The lefthand side is a piecewise-linear increasing function of  $1/\nu^*$ , with breakpoints at  $\alpha_i$ , so the equation has a unique solution which is readily determined.

This solution method is called water-filling for the following reason. We think of  $\alpha_i$  as the ground level above patch i, and then flood the region with water to a depth  $1/\nu$ . The total amount of water used is then  $\sum_{i=1}^{n} \max\{0, 1/\nu^* - \alpha - i\}$ . We then increase the flood level until we have used a total amount of water equal to one. The depth of water above patch i is then the optimal value  $x_i^*$ .

**Problem 6:** Get the optimal value of the following problem via KKT conditions:

$$\min_{x \in \mathbb{R}, y > 0} e^{-x}$$
s.t. 
$$\frac{x^2}{y} \le 0$$
(10)

**Solution:** According to the constraints, it is easy to convert the origin problem as

$$\min_{x \in \mathbb{R}} e^{-x} 
\text{s.t. } x = 0$$
(11)

The Lagrange function is then

$$L(x,\nu) = e^{-x} + \nu x$$

And the dual function is

$$g(\nu) = \inf_{x} (e^{-x} + \nu x) = \begin{cases} \nu - \nu \ln \nu & \nu > 0, \\ 0, & \nu = 0, \\ -\infty & otherwise. \end{cases}$$

Therefore, the dual problem is

$$\max_{\nu} \left\{ \begin{array}{ll} \nu - \nu \ln \nu & \nu > 0, \\ 0, & \nu = 0, \end{array} \right.$$

The optimal solution is achieved when  $\nu = 1$ . The duality gap is 0.

Problem 7: Get the dual problem of

min 
$$x^T W x$$
  
s.t.  $x_i^2 = 1, i = 1, ..., n$  (12)

**Solution:** The Lagrange function is

$$L(x,\nu) = x^T W x + \sum_{i=1}^n \nu_i (x_i^2 - 1) = x^T (W + \text{diag}(\nu)) x - 1^T \nu$$

The dual function is

$$g(\nu) = \inf_{x} x^{T} (W + \operatorname{diag}(\nu)) x - 1^{T} \nu$$

$$= \begin{cases} -1^{T} \nu & W + \operatorname{diag}(\nu) \succeq 0, \\ -\infty & otherwise. \end{cases}$$
(13)

where we use the fact that the infimum of a quadratic form is either zero (if the form is positive semidefinite) or  $-\inf$  (if the form is not positive semidefinite).

This dual function provides lower bounds on the optimal value. For example, we can take the specific value of the dual variable  $\nu = -\lambda_{\min}(W)1$ , which is dual feasible, since  $W + \operatorname{diag}(\nu) = W - \lambda_{\min}(W)I \succeq 0$ . This yields the bound on the optimal value  $p^*$ :  $p^* \geq -1^T \nu = n\lambda_{\min}(W)$ .

Problem 8: Get the conjugate function of

$$\min_{x \in \mathcal{L}} f(x) \\
\text{s.t. } x = 0$$
(14)

**Solution:** This problem has Lagrangian  $L(x,\nu) = f(x) + \nu^T x$ , and dual function

$$g(\nu) = \inf_{x} (f(x) + \nu^{T} x) = -\sup_{x} ((-\nu)^{T} x - f(x)) = -f^{*}(-\nu).$$

**Problem 9:** Get the dual problem of

$$\min_{x \in \mathcal{S}} f_0(x) 
\text{s.t. } Ax \le b, Cx = d$$
(15)

**Solution:** Using the conjugate of  $f_0$  we can write the dual function for the problem as

$$g(\lambda, \nu) = \inf_{x} (f_0(x) + \lambda^T (Ax - b) + \nu^T (Cx - d))$$

$$= -b^T \lambda - d^T \nu + \inf_{x} (f_0(x) + (A^T \lambda + C^T \nu)^T x)$$

$$= -b^T \lambda - d^T \nu - f_0^* (-A^T \lambda - C^T \nu).$$
(16)

The domain of g follows from the domain of  $f_0^*$ :

$$\operatorname{dom} g = \{(\lambda, \nu) | -A^T \lambda - C^T \nu \in \operatorname{dom} f_0^* \}.$$